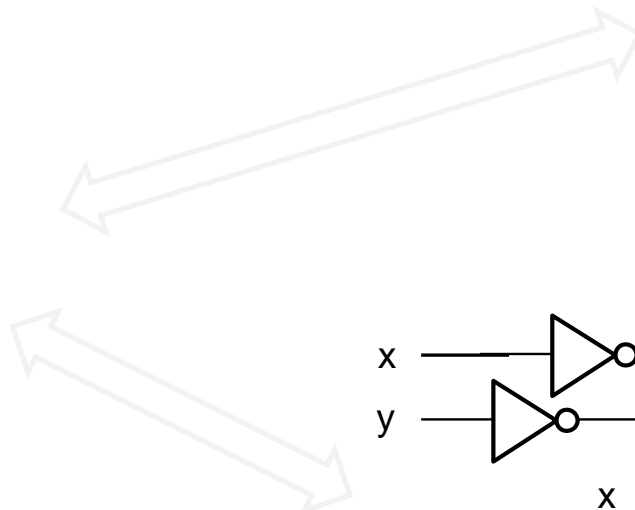


# Canonical Logic Forms

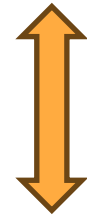
CMSC 313  
Raphael Elspas

# Converting between expression and circuit

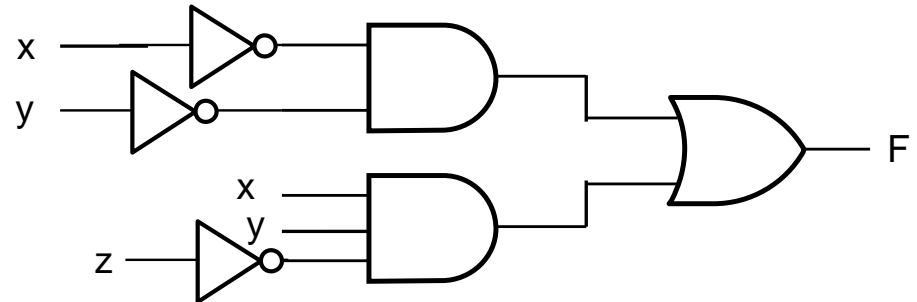
x	y	z	F
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	0



$$F = \bar{x}\bar{y} + xy\bar{z}$$

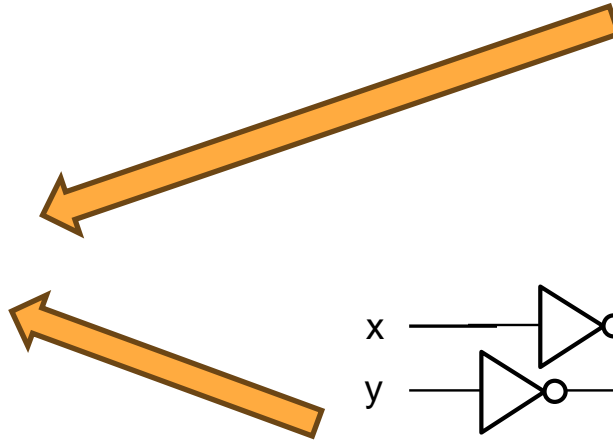


Translation

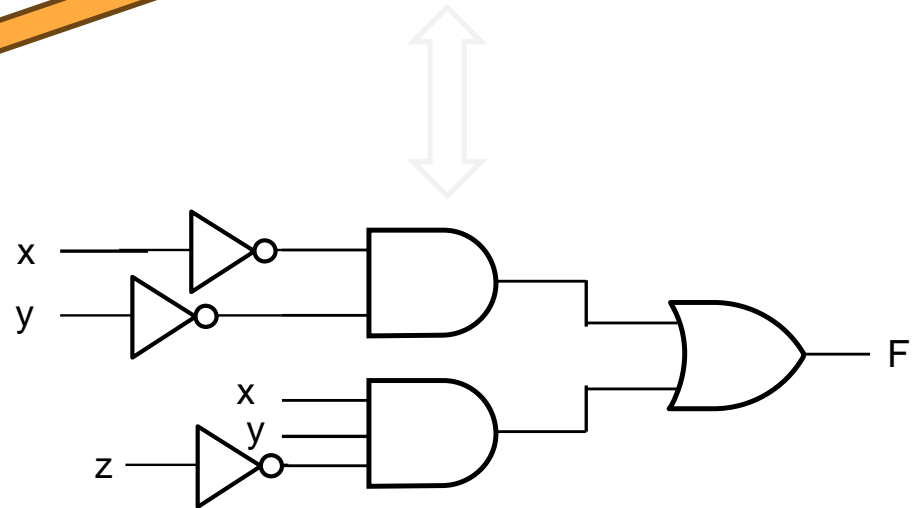


# Converting to truth table: **plug in values!**

x	y	z	F
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	0

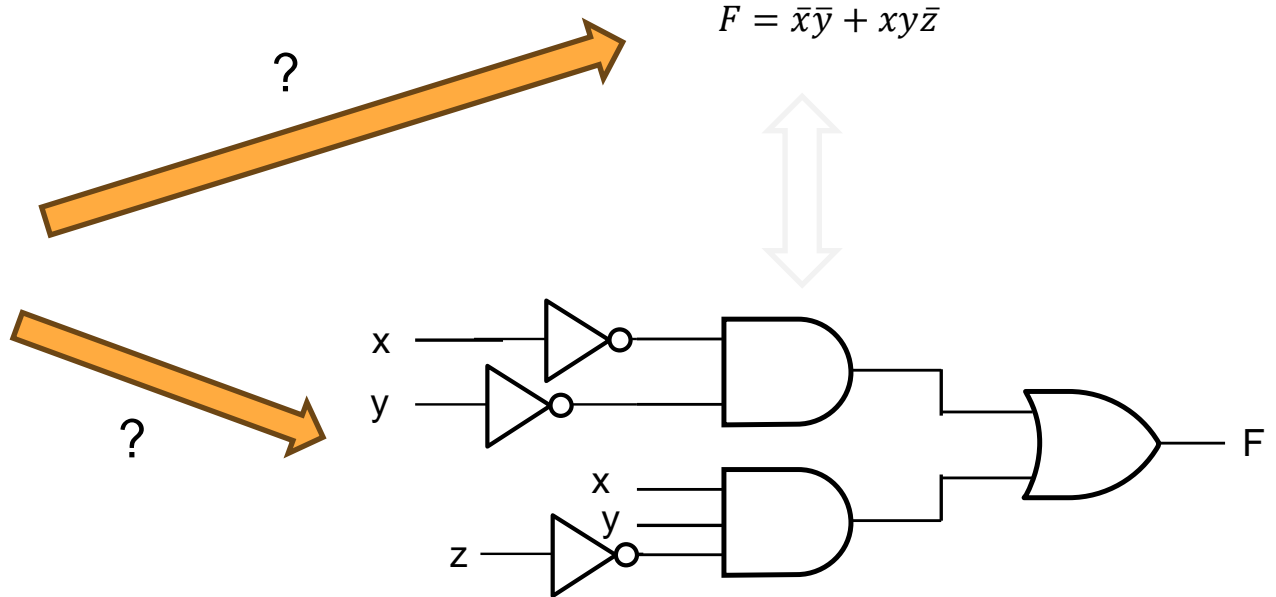


$$F = \bar{x}\bar{y} + xy\bar{z}$$



# How to convert the other way?

x	y	z	F
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	0



# Canonical Forms

- Boolean expressions that have a consistent form
- Each expression has a one to one correlation to a truth table
- Two kinds:
  - Sum of **minterms**, Sum of Products (SOP)
  - Product of **Maxterms**, Product of Sums (POS)
- Circuits that are shallower: logic has to pass through fewer circuits from input to output. This is faster because of gate delay.

# minterm

- A **minterm**, denoted as  $m_i$ , where  $0 \leq i < 2^n$ , is a product (AND) of the  $n$  variables in which each variable is
  - **complemented** if the value assigned to it is **1**, and
  - **uncomplemented** if it is **0**.
- $m_i$  is associated with the  **$i$ th row** out of  $n$  rows in the truth table
- Any Boolean function can be expressed as a sum (OR) of its minterms.
- A sum of minterms is called Sum of Products (SOP)

# minterms of 3 variables

- A shorthand notation:  
 $F(\text{list of variables}) = \Sigma(\text{list of 1-minterm indices})$

- Example:

$$\begin{aligned}
 F &= \bar{x}yz + x\bar{y}z + xy\bar{z} + xyz \\
 &= m_3 + m_5 + m_6 + m_7 \\
 &= \sum (3,5,6,7)
 \end{aligned}$$

x	y	z	minterm	notation
0	0	0	$\bar{x}\bar{y}\bar{z}$	$m_0$
0	0	1	$\bar{x}\bar{y}z$	$m_1$
0	1	0	$\bar{x}y\bar{z}$	$m_2$
0	1	1	$\bar{x}yz$	$m_3$
1	0	0	$x\bar{y}\bar{z}$	$m_4$
1	0	1	$x\bar{y}z$	$m_5$
1	1	0	$xy\bar{z}$	$m_6$
1	1	1	$xyz$	$m_7$

# Inverse of minterm

- The inverse of a sum of minterms is a sum of all the remaining minterms
- DeMorgan's application can be complicated
- Example: find inverse  $\bar{F}$

$$\begin{aligned}
 F &= \bar{x}yz + x\bar{y}z + xy\bar{z} + xyz \\
 &= m_3 + m_5 + m_6 + m_7 \\
 &= \sum (3,5,6,7)
 \end{aligned}$$

$$\begin{aligned}
 \bar{F} &= m_0 + m_1 + m_2 + m_4 \\
 &= \sum (0,1,2,4)
 \end{aligned}$$

x	y	z	minterm	F	F'
0	0	0	$\bar{x}\bar{y}\bar{z} = m_0$	0	1
0	0	1	$\bar{x}\bar{y}z = m_1$	0	1
0	1	0	$\bar{x}y\bar{z} = m_2$	0	1
0	1	1	$\bar{x}yz = m_3$	1	0
1	0	0	$x\bar{y}\bar{z} = m_4$	0	1
1	0	1	$x\bar{y}z = m_5$	1	0
1	1	0	$xy\bar{z} = m_6$	1	0
1	1	1	$xyz = m_7$	1	0



# Convert expression into SOP using a truth table

Example:  $F = x + yz$

Step 2. Derive SOP

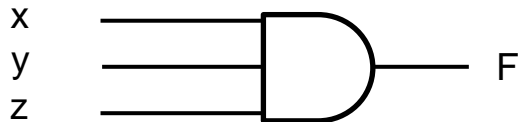
Step 1. Derive truth table

x	y	z	F
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

$$F = m_3 + m_4 + m_5 + m_6 + m_7$$
$$= \sum (3,4,5,6,7)$$

# Multiple Input Gates

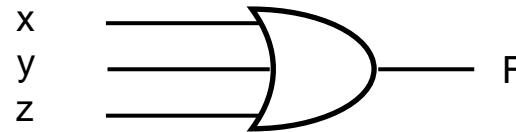
- For AND gates with input set  $S$ , if all elements in  $S$  equal 1, then output is 1. Otherwise the output is 0.



3-in AND

X	Y	Z	F
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

- For OR gates with input set  $S$ , if all elements in  $S$  equal 0, then output is 0. Otherwise the output is 1.



3-in OR

X	Y	Z	F
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

# Convert SOP into circuit

Example:  $F = m_2 + m_3 + m_7$  (for 3 variables)

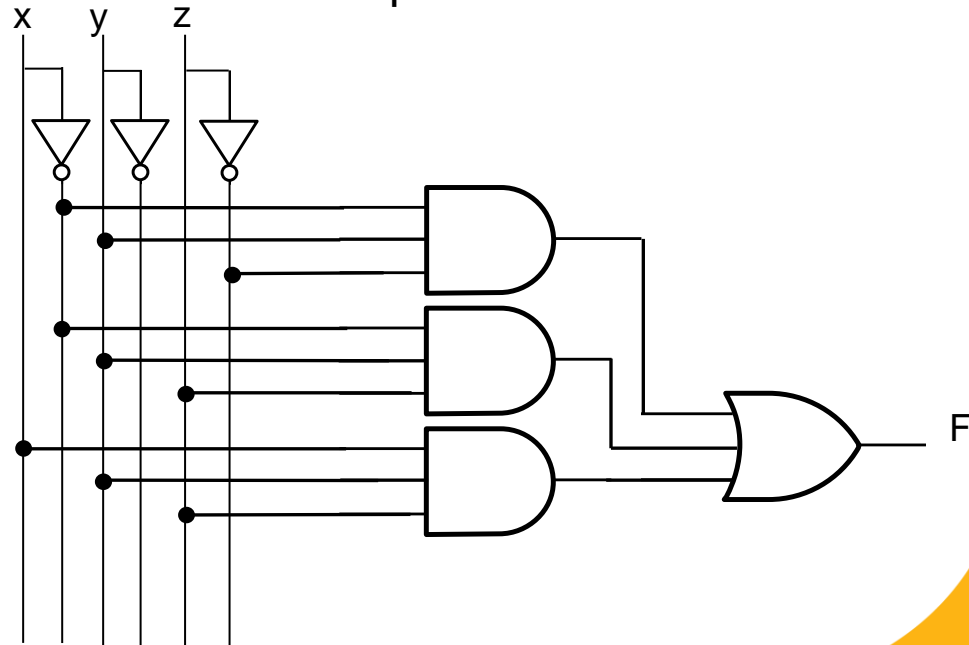
Step 1. Derive truth table

x	y	z	F
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

Step 2. Extract minterms

$$F = \bar{x}y\bar{z} + \bar{x}yz + xyz$$

Step 3. Create circuit



# Maxterm

- A **Maxterm**, denoted as  $M_i$ , where  $0 \leq i < 2^n$ , is a sum (OR) of the  $n$  variables (literals) in which each variable is
  - **complemented** if the value assigned to it is **1**, and
  - **uncomplemented** if it is **0**.
  - Note this is **reverse** of the definition for minterms
- Any Boolean function can be expressed as a product (AND) of its Maxterms.
- A product of Maxterms is called Products of Sums (POS)

# Maxterms of 3 variables

- A shorthand notation:  
 $F(\text{list of variables}) = \prod(\text{list of Maxterm indices})$
- $\prod$  is read “product of”

- Example: find  $\prod$  notation of:

$$\begin{aligned}
 F &= (x + y + z)(x + y + \bar{z})(x + \bar{y} + z)(\bar{x} + y + z) \\
 &= M_0 M_1 M_2 M_4 \\
 &= \prod(0,1,2,4)
 \end{aligned}$$

x	y	z	Maxterm	notation
0	0	0	$x + y + z$	$M_0$
0	0	1	$x + y + \bar{z}$	$M_1$
0	1	0	$x + \bar{y} + z$	$M_2$
0	1	1	$x + \bar{y} + \bar{z}$	$M_3$
1	0	0	$\bar{x} + y + z$	$M_4$
1	0	1	$\bar{x} + y + \bar{z}$	$M_5$
1	1	0	$\bar{x} + \bar{y} + z$	$M_6$
1	1	1	$\bar{x} + \bar{y} + \bar{z}$	$M_7$

# Maxterms of the **zeros** are the output!

$$\begin{aligned}
 F &= (x + y + z)(x + y + \bar{z})(x + \bar{y} + z)(\bar{x} + y + z) \\
 &= M_0 M_1 M_2 M_4 \\
 &= \prod (0, 1, 2, 4)
 \end{aligned}$$

x	y	z	Maxterm	notation	F
0	0	0	$x + y + z$	$M_0$	0
0	0	1	$x + y + \bar{z}$	$M_1$	0
0	1	0	$x + \bar{y} + z$	$M_2$	0
0	1	1	$x + \bar{y} + \bar{z}$	$M_3$	1
1	0	0	$\bar{x} + y + z$	$M_4$	0
1	0	1	$\bar{x} + y + \bar{z}$	$M_5$	1
1	1	0	$\bar{x} + \bar{y} + z$	$M_6$	1
1	1	1	$\bar{x} + \bar{y} + \bar{z}$	$M_7$	1

# Inverse of Maxterm

- The inverse of a product of Maxterms is a product of all the remaining Maxterms
- Example: Find  $\bar{F}$

$$\begin{aligned}
 F &= (x + y + z)(x + y + \bar{z})(x + \bar{y} + z)(\bar{x} + y + z) \\
 &= M_0 M_1 M_2 M_4 \\
 &= \prod (0, 1, 2, 4)
 \end{aligned}$$

$$\begin{aligned}
 \bar{F} &= M_3 M_5 M_6 M_7 \\
 &= \prod (3, 5, 6, 7)
 \end{aligned}$$

x	y	z	Maxterm	F	F'
0	0	0	$x + y + z = M_0$	0	1
0	0	1	$x + y + \bar{z} = M_1$	0	1
0	1	0	$x + \bar{y} + z = M_2$	0	1
0	1	1	$x + \bar{y} + \bar{z} = M_3$	1	0
1	0	0	$\bar{x} + y + z = M_4$	0	1
1	0	1	$\bar{x} + y + \bar{z} = M_5$	1	0
1	1	0	$\bar{x} + \bar{y} + z = M_6$	1	0
1	1	1	$\bar{x} + \bar{y} + \bar{z} = M_7$	1	0

# Convert expression into POS using a truth table

Example:  $F = x + yz$

Step 2. Derive POS

Step 1. Derive truth table

x	y	z	F
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

$$F = M_0 M_1 M_2$$
$$= \prod (0,1,2)$$



# minterm and Maxterm

Example

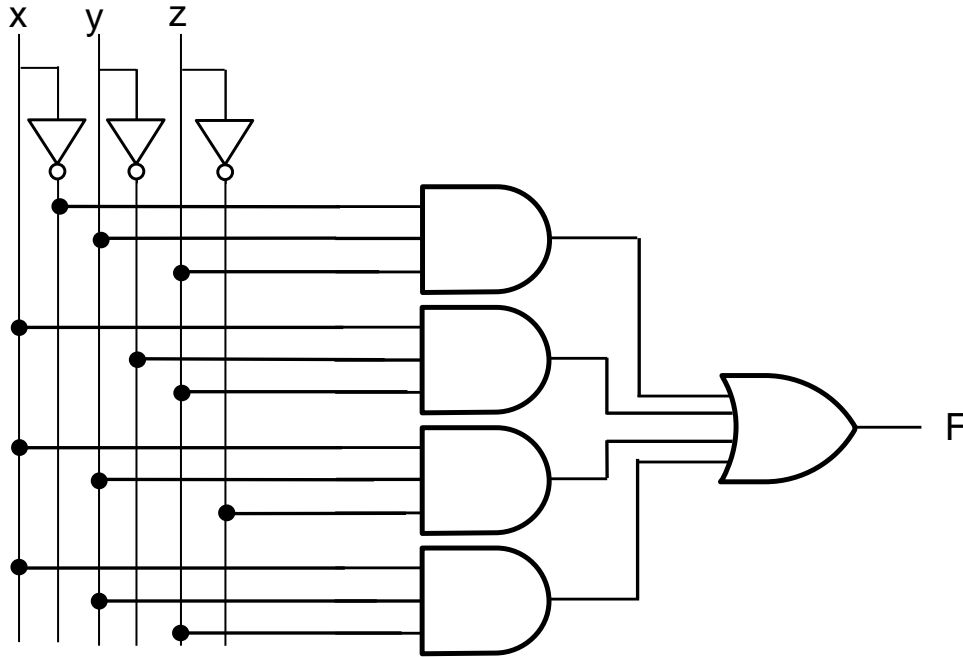
$$F \left\{ \begin{array}{l} F = \bar{x}yz + x\bar{y}z + xy\bar{z} + xyz = m_3 + m_5 + m_6 + m_7 = \sum (3,5,6,7) \\ F = (x + y + z)(x + y + \bar{z})(x + \bar{y} + z)(\bar{x} + y + z) = M_0M_1M_2M_4 = \prod (0,1,2,4) \end{array} \right.$$

$$\bar{F} \left\{ \begin{array}{l} \bar{F} = \bar{x}\bar{y}\bar{z} + \bar{x}\bar{y}z + \bar{x}y\bar{z} + x\bar{y}\bar{z} = m_0 + m_1 + m_2 + m_4 = \sum (0,1,2,4) \\ \bar{F} = (x + \bar{y} + \bar{z})(\bar{x} + y + \bar{z})(\bar{x} + \bar{y} + z)(\bar{x} + \bar{y} + \bar{z}) = M_3M_5M_6M_7 = \prod (3,5,6,7) \end{array} \right.$$

Are these (non-canonical) POS, SOP, both, neither?

- $\bar{a}b + cd$  SOP
- $c + \bar{a}$  SOP and POS
- $(c + \bar{a})(d + \bar{a} + b)$  POS
- $(c + \bar{a})d$  POS
- $(c + \bar{a})db$  POS
- $(c + \bar{a})(db + \bar{a})$  neither

# Circuits from minterms

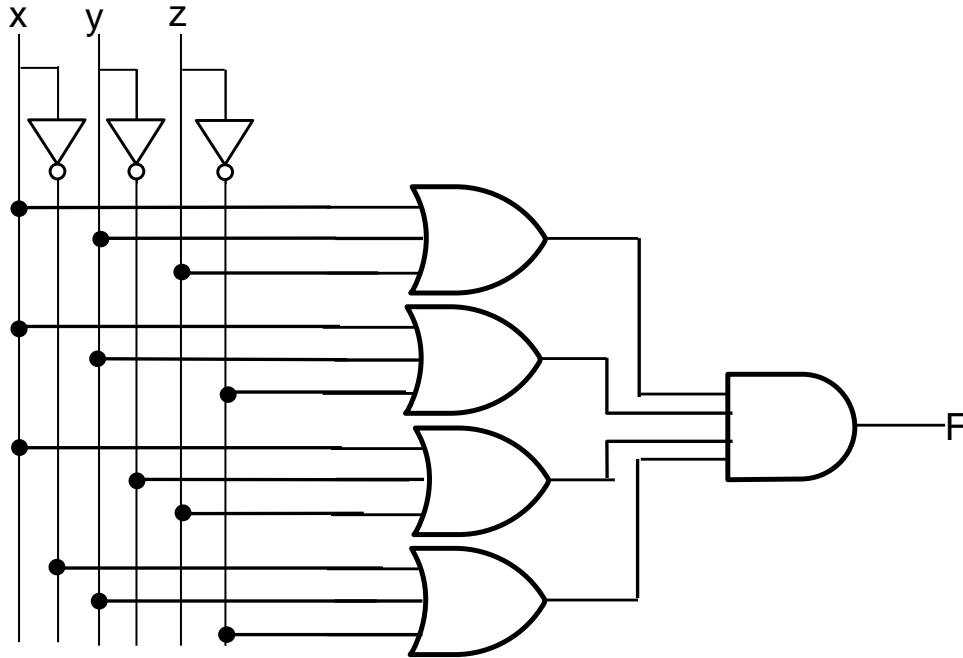


$$F = m_3 + m_5 + m_6 + m_7$$

$$= \sum (3,5,6,7)$$

x	y	z	minterm	F	F'
0	0	0	$\bar{x}\bar{y}\bar{z} = m_0$	0	1
0	0	1	$\bar{x}\bar{y}z = m_1$	0	1
0	1	0	$\bar{x}y\bar{z} = m_2$	0	1
0	1	1	$\bar{x}yz = m_3$	1	0
1	0	0	$x\bar{y}\bar{z} = m_4$	0	1
1	0	1	$x\bar{y}z = m_5$	1	0
1	1	0	$xy\bar{z} = m_6$	1	0
1	1	1	$xyz = m_7$	1	0

# Circuits from Maxterms

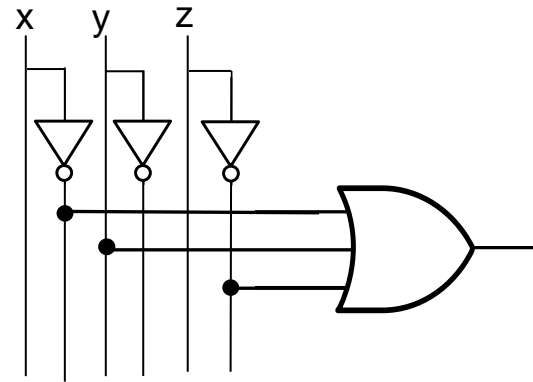
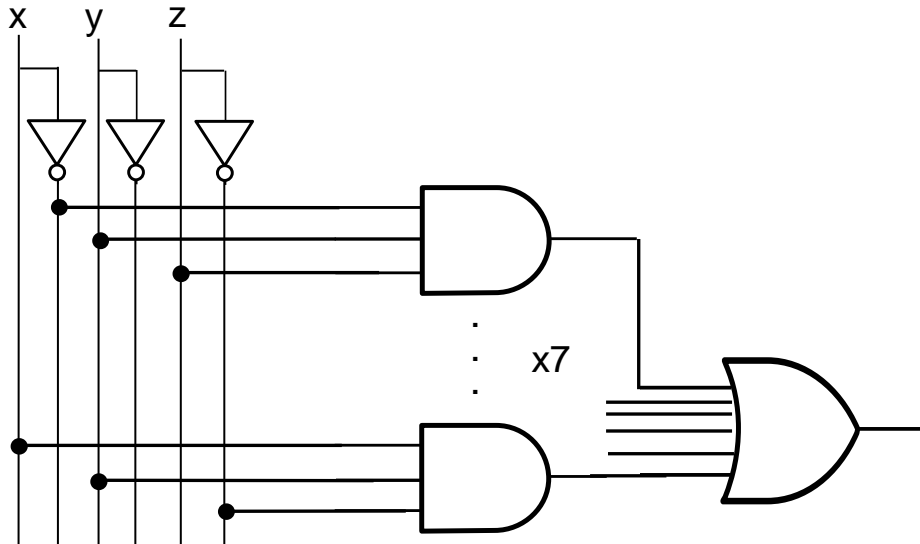


x	y	z	Maxterm	F	F'
0	0	0	$x + y + z = M_0$	0	1
0	0	1	$x + y + \bar{z} = M_1$	0	1
0	1	0	$x + \bar{y} + z = M_2$	0	1
0	1	1	$x + \bar{y} + \bar{z} = M_3$	1	0
1	0	0	$\bar{x} + y + z = M_4$	0	1
1	0	1	$\bar{x} + y + \bar{z} = M_5$	1	0
1	1	0	$\bar{x} + \bar{y} + z = M_6$	1	0
1	1	1	$\bar{x} + \bar{y} + \bar{z} = M_7$	1	0

$$\begin{aligned}
 F &= M_0 M_1 M_2 M_4 \\
 &= \prod (0,1,2,4)
 \end{aligned}$$

# minterms vs Maxterms

- For  $n$  variables, if POS uses  $x$  terms, SOP will use  $2^n - x$  terms.
- Tradeoff means simpler circuit, cheaper to manufacture
- Example:  $\Sigma(0,1,2,3,4,6,7) = \Pi(5)$ . POS is much simpler.



# Self Duals

- Reminder: a Boolean expression is **self dual** if it equals its dual. A dual is produced by replacing all ANDs with ORs and vice versa and 1s with 0s.
- New definition: A Boolean expression is self dual if:
  1. The expression is **neutral**, i.e. the number of minterms equals the number of Maxterms, and
  2. The expression does **not** contain two **mutually exclusive** terms, e.g.  $xyz$  and  $\bar{x}\bar{y}\bar{z}$  are mutually exclusive because all the variables in one term are complemented in the other.  $x\bar{y}z$  and  $\bar{x}y\bar{z}$  are also mutually exclusive.

# Self dual example

- Is  $a \oplus b$  self dual?
  1. Is the expression neutral?

a	b	$a \oplus b$
0	0	0
0	1	1
1	0	1
1	1	0

Yes, 2 minterms, 2 Maxterms



2. The expression contains **mutually exclusive** terms:  $a \oplus b = a\bar{b} + \bar{a}b$ , minterms  $m_1$  and  $m_2$  are mutually exclusive since the variables in  $a\bar{b}$  are complemented in the other:  $\bar{a}b$ .

**Not self dual**

# Self dual example

- Is  $F = \sum(3,5,6,7)$  self-dual?

x	y	z	F
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

F is an inverted mirror. Therefore, there are an equal number of minterms and maxterms, and no terms are mutually exclusive. Therefore, F is **self dual**.



## Simplest form

- Is either a minterm SOP or a maxterm POS the expression with the fewest literals? The simplest expression?
- No!
- Karnaugh maps are used to find the simplest expression and therefore a minimal literals and gates

## Summary

- Canonical Form used to convert truth table to consistent expression
- Sum of **minterms**, Sum of Products (SOP)
- Product of **Maxterms**, Product of Sums (POS)
- SOP and POS have inverse quantity of terms

# References

- <https://www.cs.ucr.edu/~ehwang/courses/cs120a/00winter/minterms.pdf>